COV886 Special Module in Algorithms: Computational Social Choice

Lecture 4

Structured Preferences

Feb 05, 2022 | Rohit Vaish

Reminder about starting recording



No group-level transitivity

Condorcet paradox





No group-level No "reasonable" transitivity voting rules

Condorcet paradox Gibbard-Satterthwaite and Arrow's theorems







No group-level transitivity No "reasonable" voting rules

Manipulation can be easy

Condorcet paradox

Gibbard-Satterthwaite and Arrow's theorems

Greedy strategy









No group-levelNo "reasonable"ManipulationNP-hardnesstransitivityvoting rulescan be easyas a barrierto manipulationCondorcetGibbard-SatterthwaiteGreedy strategyBartholdi-Tovey-Trick









No group-level
transitivityNo "reasonable"
voting rulesManipulation
can be easyNP-hardness
as a barrier
to manipulationCondorcet
paradoxGibbard-Satterthwaite
and Arrow's theoremsGreedy strategyBartholdi-Tovey-Trick

There is an ordering "<" over candidates.

a < b < c < d < e

There is an ordering "<" over candidates.

a < b < c < d < e

There is an ordering "<" over candidates.

a < b < c < d < e

A preference profile is "single-peaked" w.r.t. the ordering "<" if for each voter v:

There is an ordering "<" over candidates.

a < b < c < d < e

A preference profile is "single-peaked" w.r.t. the ordering "<" if for each voter v:

• if $top(v) \le x < y$, then v prefers x over y

There is an ordering "<" over candidates.

a < b < c < d < e

A preference profile is "single-peaked" w.r.t. the ordering "<" if for each voter v:

- if $top(v) \le x < y$, then v prefers x over y
- if $y < x \le top(v)$, then v prefers x over y

There is an ordering "<" over candidates.

a < b < c < d < e

A preference profile is "single-peaked" w.r.t. the ordering "<" if for each voter v:

- if $top(v) \le x < y$, then v prefers x over y
- if $y < x \le top(v)$, then v prefers x over y



There is an ordering "<" over candidates.

a < b < c < d < e

A preference profile is "single-peaked" w.r.t. the ordering "<" if for each voter v:

- if $top(v) \le x < y$, then v prefers x over y
- if $y < x \le top(v)$, then v prefers x over y



Single-Peaked Preferences

There is an ordering "<" over candidates. a < b < c < d < e

A preference profile is "single-peaked" w.r.t. the ordering "<" if for each voter v:

- if $top(v) \le x < y$, then v prefers x over y
- if $y < x \le top(v)$, then v prefers x over y



Single-Peaked Preferences

[Black'48]









17 18 19 20 21



17 18 19 20 21



5% 10% 15% 20%



17 18 19 20 21



5% 10% 15% 20%







17 18 19 20 21







5% 10% 15% 20%





Order the voters according to their top choices

Order the voters according to their top choices



Order the voters according to their top choices



If no. of voters is odd (say, 2k+1)

Order the voters according to their top choices



If no. of voters is odd (say, 2k+1)

Order the voters according to their top choices



If no. of voters is odd (say, 2k+1)

Order the voters according to their top choices



If no. of voters is odd (say, 2k+1)

Order the voters according to their top choices



If no. of voters is odd (say, 2k+1) If no. of voters is even (say, 2k)

Order the voters according to their top choices



If no. of voters is odd (say, 2k+1) If no. of voters is even (say, 2k)

Order the voters according to their top choices



If no. of voters is odd (say, 2k+1) If no. of voters is even (say, 2k)

Order the voters according to their top choices



If no. of voters is odd (say, 2k+1)

If no. of voters is even (say, 2k)

 $top(v_k)$ is a Condorcet winner

All candidates between top(v_k) and top(v_{k+1}) are weak Condorcet winners
Median voter rule

- 1. Each voter reports its favorite candidate (or "peak").
- 2. The* median of the reported peaks is the winner.

Median voter rule

- 1. Each voter reports its favorite candidate (or "peak").
- 2. The* median of the reported peaks is the winner.



Median voter rule

- 1. Each voter reports its favorite candidate (or "peak").
- 2. The* median of the reported peaks is the winner.



Why is the median voter rule strategyproof?



Why is the median voter rule strategyproof? Suppose the no. of voters is odd.



Why is the median voter rule strategyproof?

Suppose the no. of voters is odd.



Why is the median voter rule strategyproof?

Suppose the no. of voters is odd.



Why is the median voter rule strategyproof?

Suppose the no. of voters is odd.



Why is the median voter rule strategyproof?

Suppose the no. of voters is odd.



Why is the median voter rule strategyproof?

Suppose the no. of voters is odd.

For any misreport on the same side, the outcome doesn't change.

Any misreport on the other side moves the outcome away from the peak.



Why is the median voter rule strategyproof?

Suppose the no. of voters is odd.

For any misreport on the same side, the outcome doesn't change.

Any misreport on the other side moves the outcome away from the peak.



Why is the median voter rule strategyproof?

Suppose the no. of voters is odd.

For any misreport on the same side, the outcome doesn't change.

Any misreport on the other side moves the outcome away from the peak.



Why is the median voter rule strategyproof? Suppose the no. of voters is odd.



Why is the median voter rule strategyproof? Suppose the no. of voters is even.



Why is the median voter rule strategyproof? Suppose the no. of voters is even.



Why is the median voter rule strategyproof? Suppose the no. of voters is even.

Deterministically pick a fixed median (either left or right).



Single-peaked preferences admit a strategyproof voting rule. Why is the median voter rule strategyproof? Suppose the no. of voters is even. Deterministically pick a fixed median (either left or right). In fact, any kth order statistic works [Moulin, 1980].



Single-peaked preferences admit X strategyproof voting rules Why is the median voter rule strategyproof? Suppose the no. of voters is even. Deterministically pick a fixed median (either left or right). In fact, any kth order statistic works [Moulin, 1980].



Ok, single-peaked preferences are great. But how to find out if a given set of preferences are single-peaked?



a b c d e

b

С

C

a

е



C

a

е







e

С

a



















For this, let us discuss an equivalent definition of single-peaked preferences.

Contiguous Segments Property
A preference profile satisfies contiguous segments property w.r.t. < if, for each vote and for every k, the set of top-k candidates in that vote forms a contiguous segment w.r.t. <.

a b c d e





































A preference profile satisfies contiguous segments property w.r.t. < if, for each vote and for every k, the set of top-k candidates in that vote forms a contiguous segment w.r.t. <.



Contiguous segments (w.r.t. <) ⇔ Single-peaked (w.r.t. <)

Contiguous Segments ⇒ Single-Peaked

Suppose the contiguous segments property holds w.r.t. the axis <.

Suppose the contiguous segments property holds w.r.t. the axis <.

Suppose the contiguous segments property holds w.r.t. the axis <.



Suppose the contiguous segments property holds w.r.t. the axis <.



Suppose the contiguous segments property holds w.r.t. the axis <.



Suppose the contiguous segments property holds w.r.t. the axis <.



Suppose the contiguous segments property holds w.r.t. the axis <.

Suppose, for contradiction, that for some pair of candidates x,y and some voter v, top(v) < x < y but v prefers y over x.



x disconnects the top-k segment

top(v)

$$y \longrightarrow k^{th}$$
 ranked
 \vdots
 $x \longrightarrow (k+1)^{th}$ or below

Contiguous Segments ⇒ Single-Peaked

Suppose the contiguous segments property holds w.r.t. the axis <.

Suppose, for contradiction, that for some pair of candidates x,y and some voter v, top(v) < x < y but v prefers y over x.



x disconnects the top-k segment

top(v)

$$y \longrightarrow k^{th}$$
 ranked
 \vdots
 $x \longrightarrow (k+1)^{th}$ or below

Suppose single-peaked property holds w.r.t. the axis <.

Suppose single-peaked property holds w.r.t. the axis <.

Then, for any x < y < z, a voter v will never rank y below x and z. ("no valleys" property)

Suppose single-peaked property holds w.r.t. the axis <.

Then, for any x < y < z, a voter v will never rank y below x and z. ("no valleys" property)



Suppose single-peaked property holds w.r.t. the axis <.

Then, for any x < y < z, a voter v will never rank y below x and z. ("no valleys" property)

if top(v) is to the left of y then y is preferred over z



Suppose single-peaked property holds w.r.t. the axis <.

Then, for any x < y < z, a voter v will never rank y below x and z. ("no valleys" property)

if top(v) is to the left of y then y is preferred over z if top(v) is to the right of y then y is preferred over x


Suppose single-peaked property holds w.r.t. the axis <.

Then, for any x < y < z, a voter v will never rank y below x and z. ("no valleys" property)

if top(v) is to the left of y then y is preferred over z if top(v) is to the right of y then y is preferred over x



Thus, single peaked (w.r.t. <) \Rightarrow no valleys (w.r.t. <).

Single-Peaked \Rightarrow Contiguous Segments

Suppose single-peaked property holds w.r.t. the axis <.

Then, for any x < y < z, a voter v will never rank y below x and z. ("no valleys" property)



Suppose, for contradiction, that for some voter v, the set of top k candidates is not contiguous (w.r.t. <).

Suppose, for contradiction, that for some voter v, the set of top k candidates is not contiguous (w.r.t. <).

Pick the smallest such k.

Suppose, for contradiction, that for some voter v, the set of top k candidates is not contiguous (w.r.t. <).

Pick the smallest such k.



Suppose, for contradiction, that for some voter v, the set of top k candidates is not contiguous (w.r.t. <).

Pick the smallest such k.



Suppose, for contradiction, that for some voter v, the set of top k candidates is not contiguous (w.r.t. <).

Pick the smallest such k.



Let y be a candidate that separates x from the top (k-1) candidates

Suppose, for contradiction, that for some voter v, the set of top k candidates is not contiguous (w.r.t. <).

Pick the smallest such k.



Let y be a candidate that separates x from the top (k-1) candidates

Then, top(v), y and x constitute a valley.

Suppose, for contradiction, that for some voter v, the set of top k candidates is not contiguous (w.r.t. <).

Pick the smallest such k.



Let y be a candidate that separates x from the top (k-1) candidates

Then, top(v), y and x constitute a valley.

We will use the contiguous segments property to design an algorithm for recognizing single-peaked preferences.

We will use the contiguous segments property to design an algorithm for recognizing single-peaked preferences.

But, before that, another digression.





Given a 0-1 matrix, is there a permutation of columns such that all 1's in each row appear consecutively?



[Booth and Leuker, JCSS 1976]

The consecutive 1's problem can be solved in polynomial time.

adbceadccbadebe



abcde





a b c d e 1 0 1 0 0

abcde

1

1

0 0

0

0

0

adbceadccbadebe

abcde

1 1 0



abcde

0

U

1 1 0

0

0

1

adbceadccbadebe

adbceadccbadebe

 1
 0
 1
 0
 0

 1
 0
 1
 1
 0

 1
 1
 1
 1
 0

abcde

abcde

adbceadccbadebe

 1
 0
 1
 0
 0

 1
 0
 1
 1
 0

 1
 1
 1
 1
 0

 1
 1
 1
 1
 0

 0
 0
 0
 1
 1



1010010110111100001100111

abcde



adbceadccbadebe

abcde

adbceadccbadebe

adbceadccbadebe

a d b c e a d c c b a d e b e

adbceadccbadebe
Recognizing Single-Peaked Prefs w.r.t. some axis a b c d e a d b

c e a d c c b a d e b e

Recognizing Single-Peaked Prefs w.r.t. some axis

a d b c e a d c c b a d e b e

[Bartholdi and Trick, ORL 1986] A preference profile is single-peaked if and only if its prefix matrix satisfies consecutive 1's property.

c d e b а () 1 n 1 0 0 0

Recognizing Single-Peaked Prefs w.r.t. some axis е b a а е a C е [Bartholdi and Trick, ORL 1986] A preference profile is single-peaked 0 if and only if

its prefix matrix satisfies consecutive 1's property.

Recognizing Single-Peaked Prefs w.r.t. some axis

b

acde h

1

0

0

0

0

()

a a е С С a d e е

[Bartholdi and Trick, ORL 1986] A preference profile is single-peaked if and only if its prefix matrix satisfies consecutive 1's property. Do single-peaked preferences occur in real world?

PrefLib: A Library for Preferences

	Data 🔻	Elections	Matchings	Ratings	Search	About
	PrefLib is a reference library of preference data and links assembled by Nicholas Mattei, Toby Walsh and lately Simon Rey. This site and library is proudly supported by the Algorithmic Decision Theory group at Data61 and the The COMSOC Group at the University of Amsterdam. We want to provide a comprehensive resource for the multiple research communities that deal with preferences, including computational social choice, recommender systems, data mining, machine learning, and combinatorial optimization, to name just a few. For more information on PrefLib and some helpful tips on using it, please see Nick's Tutorial Slides and Code from EXPLORE 2014. Check out the data type page to learn more about the kind of data we provide. Please see the about page for information about the site, contacting us, and our citation policy. We rely on the support of the community in order to grow the usefulness of this site. To contribute, please contact Nicholas Mattei at: nsmattei {at} gmail or Simon Rey at: s.j.rey{at}uva{dot}nl.					In Brief
						We currently host:11 types of data
						 38 datasets 3668 data files
						More than 3.37 Gb. of data Other Links
						Here are some links that you might find relevant as well.
						DEMOCRATIX: A Declarative Approach

to Winner Determination

CHAPTER 15

A PREFLIB.ORG Retrospective: Lessons Learned and New Directions

Nicholas Mattei and Toby Walsh

Trends in Computational Social Choice

Realism. Perhaps the key motivating factor behind assembling PREFLIB was a desire to have realistic data. Many of the models studied in classical social choice seem to be chosen because they *seem* reasonable or were explicitly chosen for mathematical expediency. Perhaps nothing is more of an exemplar here than the fact that out of over 300 profiles containing strict, complete preference relations, absolutely none are single-peaked, a common profile restriction that has been called "natural" or "well-motivated" numerous times since its introduction by Black (1948). Collecting data has helped us to quantify what is reasonable. Now we have to start using the data.



Two candidates: Majority!

Three+ candidates: Majority relation can be cyclic Many voting rules, many applications!



Two candidates: Majority!

Three+ candidates: Majority relation can be cyclic Many voting rules, many applications! Any "reasonable" voting rule is either manipulable or a dictatorship.



Two candidates: Majority!

Three+ candidates: Majority relation can be cyclic Many voting rules, many applications! Any "reasonable" voting rule is either manipulable or a dictatorship.



Computational complexity can sometimes prevent manipulation.

(and a bit about sports)



Two candidates: Majority!

Three+ candidates: Majority relation can be cyclic Many voting rules, many applications! Any "reasonable" voting rule is either manipulable or a dictatorship.



Structured preferences can help circumvent negative results.

Computational complexity can sometimes prevent manipulation.

(and a bit about sports)

Next Time

House allocation and kidney exchange





Reminder: Project groups due today!

References

• Single-peaked preferences in theory:

Duncan Black "On the Rationale of Group Decision-Making" Journal of Political Economy, Feb 1948, 56(1), pg 23-34 https://www.journals.uchicago.edu/doi/10.1086/256633

• Single-peaked preferences in the real-world:

Nicholas Mattei and Toby Walsh "A PREFLIB.ORG Retrospective: Lessons Learned and New Directions" Chapter 15 in Trends in Computational Social Choice https://research.illc.uva.nl/COST-IC1205/BookDocs/TrendsCOMSOC.pdf

References

 Lecture by Edith Elkind on restricted preference domains: <u>https://www.youtube.com/watch?v=q8vc8Znoev0</u> <u>https://www.youtube.com/watch?v=vL_U-5tlQu4</u>

• Strategyproof voting rules using "phantom" voters:

Hervé Moulin "On Strategyproofness and Single-Peakedness" Public Choice, 35(4), 1980, pp. 437-455 https://www.jstor.org/stable/30023824

References

• Recognizing single-peaked preferences:

John Bartholdi III and Michael A. Trick "Stable Matching with Preferences Derived from a Psychological Model" Operations Research Letters, 5(4), 1986, pp. 165-169 https://www.sciencedirect.com/science/article/pii/0167637786900726

• Consecutive 1's Problem:

Kellogg S. Booth and George S. Lueker "Testing for the Consecutive One Property, Interval Graphs, and Graph Planarity using PQ Tree Algorithms" JCSS, 13(3), 1976, pp. 335-379 https://www.sciencedirect.com/science/article/pii/S0022000076800451